Speedier snail signatures

Thomas Decru Lorenz Panny Frederik Vercauteren

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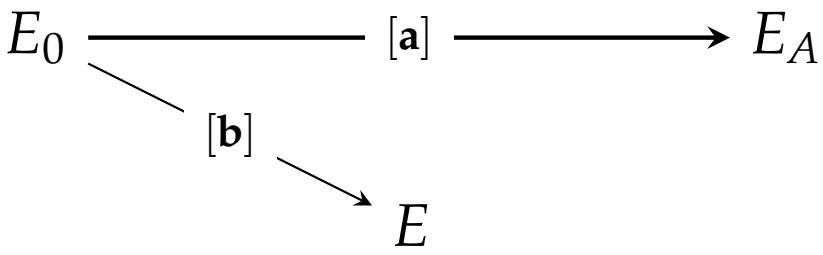

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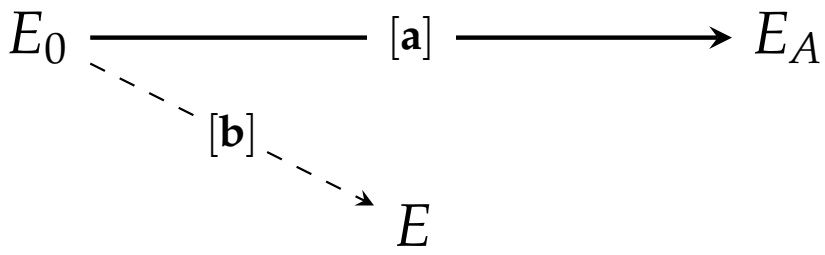

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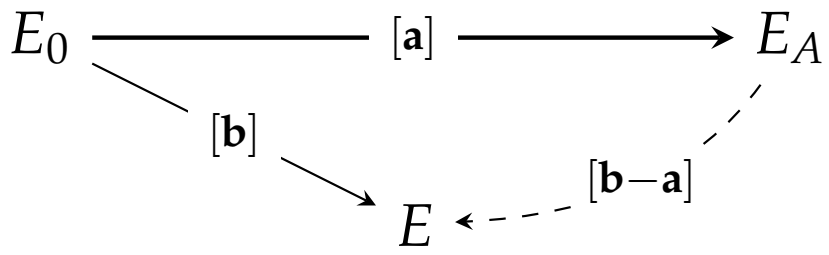

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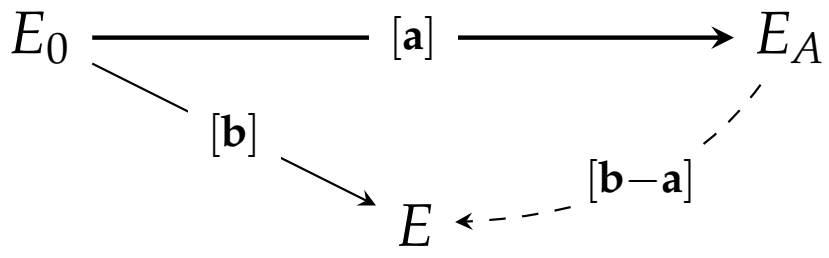

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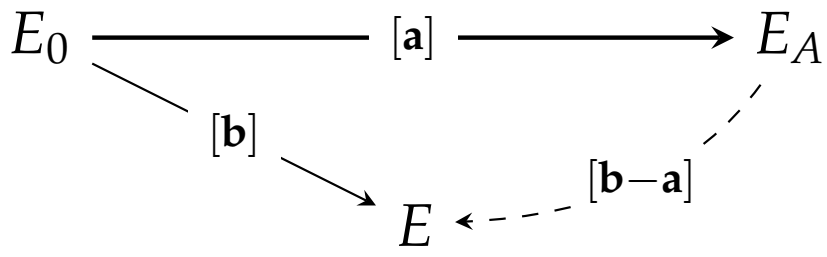

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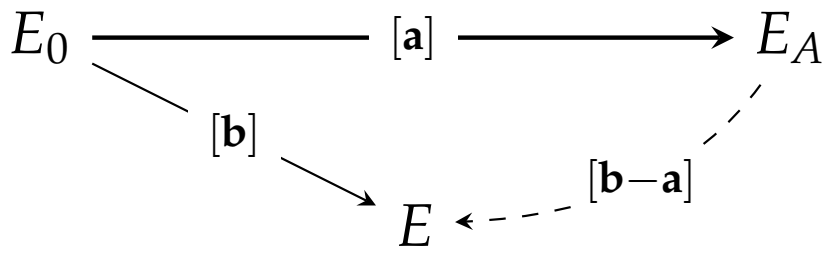

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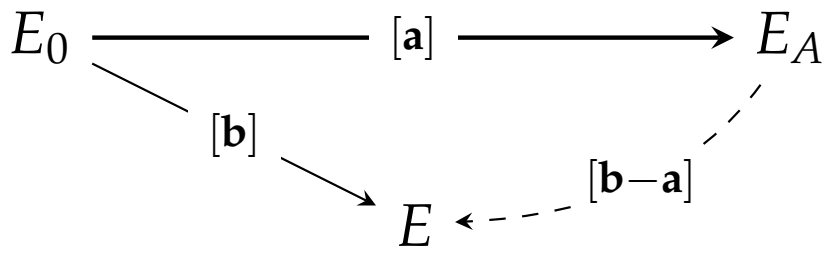

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$\rightsquigarrow$ Distribution of revealed vectors $\mathbf{b} - \mathbf{a}$ is independent of $\mathbf{a}$.

- ...but tiny rejection probabilities require huge boxes.

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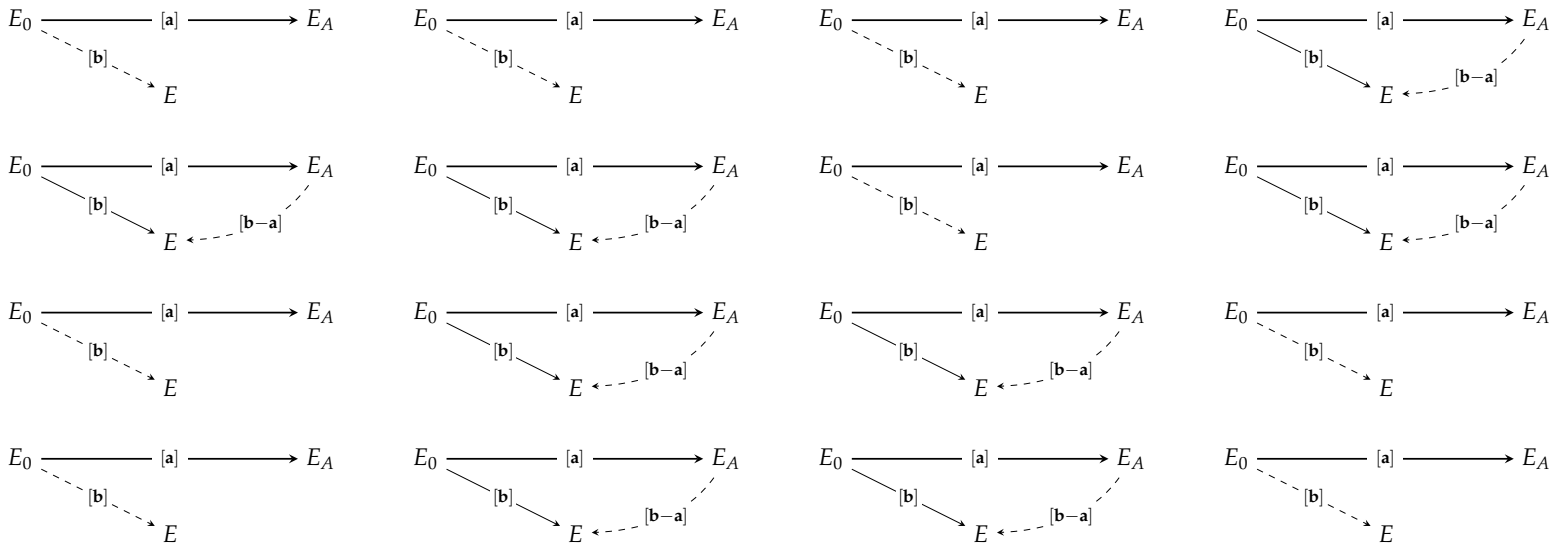

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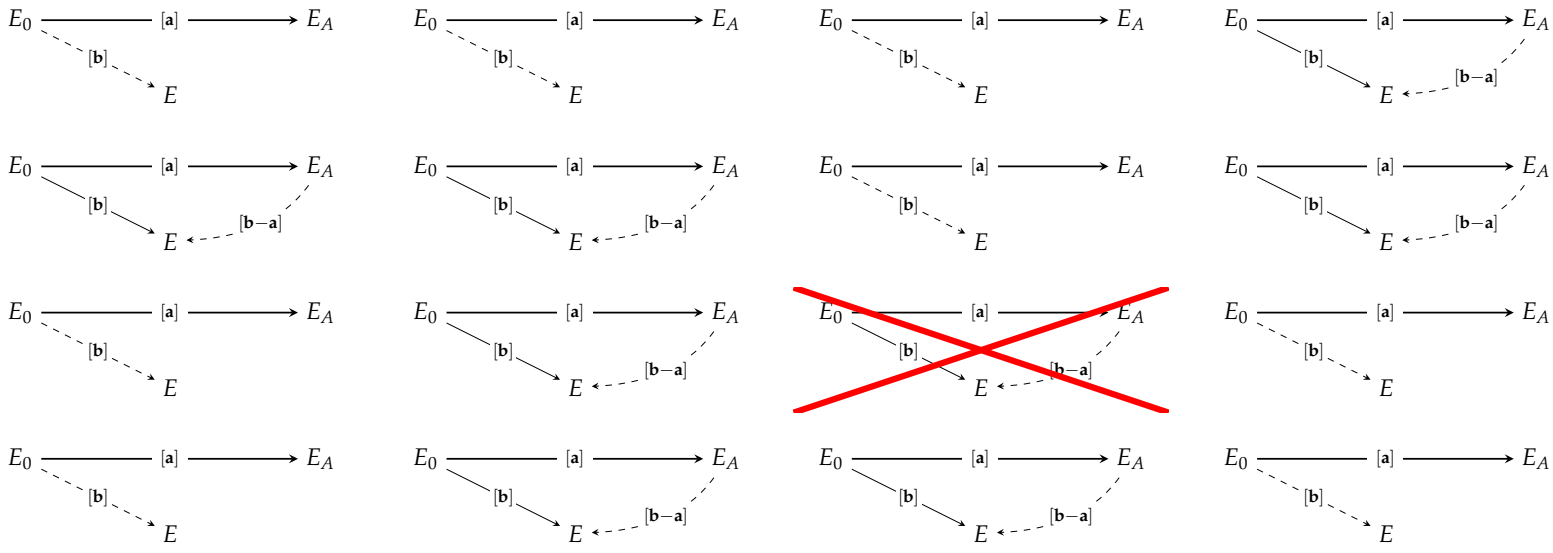

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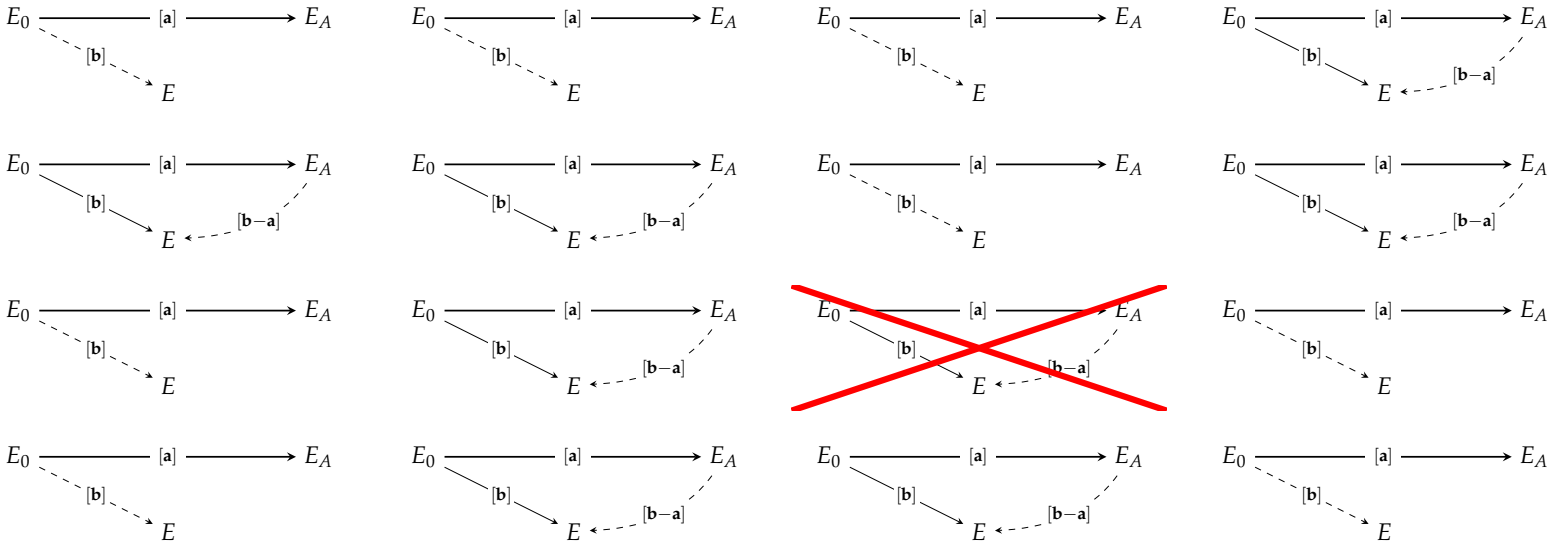

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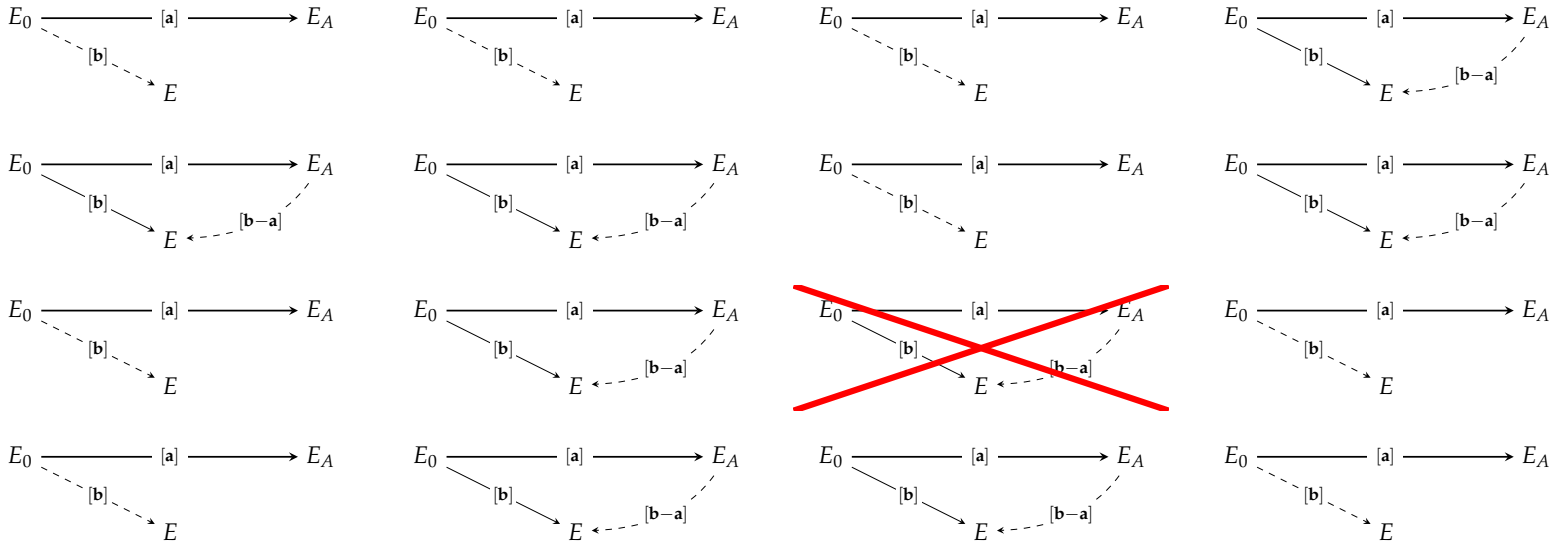

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- Our idea: Why not allow 'a few' rejections?

$\rightsquigarrow$ need $> \lambda$ correct answers, but can use smaller boxes.
$\rightsquigarrow$ better performance overall!

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